SpiderNet: Hybrid Differentiable-Evolutionary Architecture Search via Train-Free Metrics

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Abstract

Neural Architecture Search (NAS) algorithms are intended to remove the burden of manual neural network design, and have shown to be capable of designing excellent models for a variety of well-known problems. However, these algorithms require a variety of design parameters in the form of user configuration or hard-coded decisions which limit the variety of networks that can be discovered. This means that NAS algorithms do not eliminate model design tuning, they instead merely shift the burden of where that tuning needs to be applied. In this paper, we present SpiderNet, a hybrid differentiable-evolutionary and hardware-aware algorithm that rapidly and efficiently produces state-of-the-art networks. More importantly, SpiderNet is a proof-of-concept of a minimally-configured NAS algorithm; the majority of design choices seen in other algorithms are incorporated into SpiderNet’s dynamically-evolving search space, minimizing the number of user choices to just two: reduction cell count and initial channel count. SpiderNet produces models highly-competitive with the state-of-the-art, and outperforms random search in accuracy, runtime, memory size, and parameter count.

1. Introduction

Neural Architecture Search (NAS) looks to automate network design, to find quick and efficient ways of producing state-of-the-art networks for novel problems. Differentiable NAS is a popular subcategory of NAS algorithms, which makes architectural selection a continuous process and thus can be handled by gradient descent during model training, and has shown to be a tremendously efficient means of performing NAS as compared to very expensive evolutionary [12,13] or reinforcement-learning [10,17] based approaches.

While differentiable NAS algorithms can reliably produce good models, their search flexibility is implicitly limited by their configuration, that is, the parameters that define the supernet within which candidate models are isolated from. In algorithms like DARTS [6], PC-DARTS [15], or ProxylessNAS [1], the search supernet and thus the NAS search space is fixed throughout the search space, therefore constraining the available models based on the particular construction of said supernet. The design of the supernet is often either hardcoded (such as the node count and degree in DARTS) or left to user configuration (such as the cell stage counts in ProxylessNAS). In essence, this forces the search algorithm into a specific network design space as described by Radosavovic et al. [11], but Radosavovic et al. show that particular choice of design space can have highly significant effects on the performance of models within the design space population.

Given that the intended use-case of NAS is to automatically find optimal architectures for novel problems with minimal manual tuning, requiring users to define the search space is counter-productive, especially for novel problems where no network design best-practices exist. Arguably, NAS algorithms that rely on user configuration of the supernet are just moving the burden of search one level higher, replacing architecture search with architecture-search search: instead of spending time tweaking which operation goes where, time is spent tweaking things like per-cell node count or reduction cell placement.

With that question in mind, it’s worth looking at algorithms that approach NAS from the opposite direction: rather than specifying the structure of the initial supernet and using differentiable NAS to find some optimal subnet, what about algorithms that start at some minimal initial state and dynamically expand the search space as necessary? By letting the NAS algorithm specify the design of its supernet via this expansion, the flexibility and ease-of-applicability of NAS could be greatly improved.

While evolutionary NAS algorithms like Regularized Evolution [13] exactly fit this required paradigm, they are plagued by the immense computational costs required to evaluate each candidate in the gene pool. To rectify this,
we present SpiderNet\(^1\) a hybrid evolutionary-differentiable NAS algorithm that uses the train-free Neural Tangent Kernel-Linear Region Count (NTKLRC) metric [2] to rapidly evolve its supernet, while using conventional differentiable techniques to isolate promising subgraphs within. SpiderNet is a minimally-configured hardware-aware algorithm that designs highly performant networks, and notably demonstrates an overlooked domain of NAS’ advantage over random search: efficiency.

In this paper, we will first cover related work to SpiderNet in Section 2, followed by the design of SpiderNet in Section 3. Next, Section 4 describes the various experiments conducted on SpiderNet, the results of which are detailed in Section 5. Finally, sections 6 and 7 discuss these results and describe their ramifications to NAS as a whole.

2. Related Work

Evolutionary NAS was first popularized by Neuroevolution [13], which suffered from immense computational costs. Meanwhile, differentiable NAS approaches were introduced by Differentiable Architecture Search (DARTS) [6], before being refined by papers like Partially Connected DARTS (PC-DARTS) [15] and ProxylessNAS [1]. Meanwhile, train-free NAS techniques are relatively new, with papers like Training-Free NAS (TENAS) [2] and NAS Without Training (NASWOT) [8] pioneering algorithms that identify promising architectures without the need for expensive training operations like backpropagation.

3. Design

3.1. Models

The SpiderNet supernet is cell structured, where each cell contains edges and nodes. Edges consist of parallel operations, initially one of each from the 7 available in the operation search space: Identity, 3x3 Max Pool, 3x3 Average Pool, 3x3 Separable Convolutions, 5x5 Separable Convolutions, 3x3 Dilated Convolutions, and 5x5 Dilated Convolutions. These are chosen to align with the common set of computer-vision operations chosen by algorithms like DARTS [6] or ProxylessNAS [1], but any set could just as easily be chosen. However, unlike DARTS and ProxylessNAS, edges need not be singular; they can have any arbitrary mixture of operations within the operational set.

Each node performs tensor summation, to ensure that tensor shapes do not change as the number of inbound edges at a node vary. Each cell has an input node for each previous cell and the original model input; therefore, the fourth cell has four input nodes (one for Cell 1, 2, and 3, plus the model input), the third has three, etc. Each cell is initialized with just its input nodes, one intermediate node, and a single output node. One edge is then placed between each input and the intermediate node, and then a single edge is placed between the sole intermediate node and the output node. This configuration is shown in Figure 1:

This starting point can be thought of as a suitably ‘minimal’ starting point as it provides two edges per input: one to the single intermediate and one to the output. Meanwhile, the single intermediate has an edge to the output that combines and synthesizes information from the inputs. Evolutions of this architecture can therefore reinforce the pathways from the input or strengthen the ‘mixing’ portion of the cell.

On a macro-structural level, SpiderNet takes inspiration from Xie et al. [14], who use a large-cell paradigm for their model design, referring to models that use a small number of high-node cells, as opposed to the high cell-count low-node cells seen in small-cell paradigms like those used by DARTS [6] or ProxylessNAS [1]. The rationale for this is that any cellular pattern comprised of a single reduction cell followed by any number of sequential ‘normal’ cells (using the standard definition of ‘normal’ and ‘reduction’ cells introduced by Zoph et al. [18]) can be replaced as a single reduction cell of sufficiently large node count, while any repeated number of normal cells can be replaced by a single sufficiently large normal cell. For example, the DARTS model used for CIFAR10 has 20 small cells, with reduction cells at positions 7 and 14. This structure can be divided into three groups of cells \(\mathrm{rnnnnnnnnnnnnnnnnnnnnn}^2\), which by the above rules reduces to three large cells \(\mathrm{NRR}\). The advantage of this is that it eliminates all cellular design choices beyond reduction count; any possible sequence of small-cells with \(r\) reductions can be represented by \(r + 1\) large cells: one large normal cell followed by \(r\) large reduction cells.

Like in Cai et al. [1] and Xie et al. [14], SpiderNet cells are unique throughout the model; each cell has a unique node count, edge connectivity, and each edge has a unique operational mixture.

\(^1\)Code at https://github.com/RobGeada/SpiderNet-NTKLRC

\(^2\)Here \(n\) and \(r\) represent normal and reduction cells respectively.
3.2. Pruning

To identify promising subnetworks within the supernet, SpiderNet makes use of differentiable pruning, introduced by Kim et al. [4] and first used for NAS in Geada et al. [3]. Here, a weighted operation is placed along every edge we wish to make prunable, that performs the following operation over the inbound tensor $x$:

$$P(x, w) = (\text{Gate}(w) + \text{Saw}(w)) x$$  \hspace{1cm} (1)

$$\text{Gate}(w) = \begin{cases} 0 & w < 0 \\ 1 & w \geq 0 \end{cases}$$  \hspace{1cm} (2)

$$\text{Saw}(w) = \frac{Mw - \lfloor Mw \rfloor}{M}$$  \hspace{1cm} (3)

where $M$ is some large constant, in this case $10^9$. This provides a function with the mathematical properties of a binary gate, but with a constant gradient of 1 with respect to the pruner weight $w$. This allows for $w$ to be easily learned, meaning operations can be dynamically pruned during gradient descent to benefit task performance. We can then use this behavior to isolate promising subgraphs of the SpiderNet supernet.

Our usage of pruners aligns with Geada et al. [3], where differentiable pruners are placed along each parallel edge operation. As the model trains, the pruner weights adapt and switch operations on and off as necessary. Any pruner which remains switched off for more than 75% of the batches from the previous 4 epochs is deadheaded, that is, permanently deleted from the model. This ‘locks-in’ the pruner decision, freeing up memory space from the unnecessary operation which we can use for subsequent mutations of the network.

3.3. Mutation

Dynamically expanding the search space is accomplished via triangular mutations, which acts on an arbitrary edge $A \rightarrow B$ that connects two nodes $A$ and $B$ within a model, adding a third node $C$ and two new edges. These edges are wired such that for the nodes $A$, $B$, and $C$, one node sends two outputs, one node receives one input and sends one output, and one node receives two inputs, as shown in Figure 2. Each new edge is initialized with every operation from the available operation set. The triangular mutation is composable with edge pruning to create an effectively infinite variety of graph patterns, since it crucially introduces a branching path into a linear node relationship. This creates a seed for further branching by later mutations along any of these three edges, which allows a small number of mutations to create immense cellular complexity. For architectural examples of models generated

3.4. Mutation Practicalities

When mutations are performed during model training, they alter the distributions that downstream operations are expecting, worsening the model’s performance. In reaction to this, the differentiable pruners would instantly prune away the new operations, thus restoring the connectivity back to the model’s pre-mutation state. To remedy this, the model weights are reinitialized after each mutation, which allows mutation to be a productive operation.

3.5. Selecting Edges to Mutate

With the edge mutation operation defined, the next step is to identify the mutation criteria; how should the model select an edge to perform this mutation on? To perform this both rapidly and intelligently, we adapt NTK-LRC to guide mutations.

3.6. NTK-LRC Train-Free Metric

A train-free metric refers to methods that evaluate a model’s quality via sampling methods that involve no back-propagation or weight sampling, ideally in such a way that the metric correlates strongly with final model performance. Following Chen et al. [2], we use the neural tangent kernel (NTK) and linear region count (LRC) to measure a potential architecture’s trainability and expressability respectively, where trainability refers to how rapidly an architecture trains to good parameter values, while expressability describes the flexibility of the range of functions the architecture can approximate. Like Chen et al., we use the joint NTK-LRC metric which jointly represents optimal NTK and LRC via the summation of descending NTK rank to ascending LRC rank as compared to other models within the search population.

3.7. Searching via Joint NTK-LRC Metric

While in theory the NTK-LRC metric allows for every candidate in the search space to be evaluated ‘for-free’ in terms of model training, the actual computation of the two metrics still comes with a non-trivial computation cost. For
larger search spaces the majority of the space cannot be fully enumerated, and must be traversed in some intelligent manner. In the case of the DARTS space, Chen et al. uses a tournament selection [9] algorithm to find good architectures. Here, a ‘champion’ model \( M \) is compared against a set of challenger models \( \mathcal{M} \), each \( M' \in \mathcal{M} \) created by removing one operation at random from \( M \). The model \( M' \in \mathcal{M} \) with the best joint NTK-LRC is then chosen as the next champion, and the process repeats until a satisfactory model is found. In essence, this is a train-free method of iteratively pruning a supernet to find an optimal subnet, and from this approach Chen et al. find state-of-the-art results comparable to those found by DARTS.

We adapt Chen et al.’s tournament selection to measure the quality of potential mutations. Given some model \( M \), the set of all models \( \mathcal{M} \) that are one mutation away from \( M \) is enumerated. The NTK and LRC of each mutated model \( M' \in \mathcal{M} \) is compared against the original unmutated model \( M \), and the model \( M' \in \mathcal{M} \) is selected that meets the following conditions, if any such model exists:

1. Is non-detrimental to both NTK and LRC as compared to \( M \)
2. Maximizes the combined NTK-descending and LRC-ascending ranks compared to all other models in \( \mathcal{M} \)

This process can then be repeated a number of times, in theory improving the quality of the model after each cycle.

One practical caveat to this approach is the aforementioned cost of evaluating NTK and LRC; if each takes around a minute to evaluate each potential mutation, exhaustively comparing every possible mutation can take around an hour for larger models. To avoid this, mutation candidates are instead sampled at random, and after a certain configurable number of ‘good’ mutations are found (i.e., ones that do not increase NTK and do not decrease LRC, therefore producing a model that is at worst equivalent to the original model) the sampling ends. Of these found good mutations, the one with the best NTK-LRC is performed.

With this, NTK-LRC can now be used as a mutation metric, and Algorithm 1 describes how the NTK-LRC mutation metric selects candidate edges for mutation.

The NTK-LRC algorithm works by first creating a slim model \( S \) of our SpiderNet model \( M \), as described in Chen et al. This is done by initializing an exact copy of the SpiderNet model, and then setting the channel count of each operation to 1. Next, line 2 selects edges at random from the model as potential mutation candidates. Before trialling the mutation of some selected edge \( e \), line 3 creates an identical copy \( S' \) of our slim model \( S \); all trial mutations are performed on the copy, thus keeping the original slim model \( S \) as an unmodified template that can be used again. Edge \( e \) is then mutated within \( S' \) in line 4, thus creating two new edges \( e' \) and \( e'' \) as per the rules of triangular mutation. Next, a copy of the mutated \( S' \) is made in line 5. This creates two identical mutated models, the original \( S' \) and the copy, which are then labeled as \( S'_{\text{on}} \) and \( S'_{\text{off}} \). Within \( S'_{\text{off}} \), the new edges \( e' \) and \( e'' \) created by the mutation are disconnected (simply by multiplying their outputs by 0); this means that \( S'_{\text{off}} \) is now mathematically identical to the original unmutated model \( S \), but physically has the same parameters as \( S'_{\text{on}} \). See Figure 3 for a simplified visualization of \( S, S'_{\text{on}}, \) and \( S'_{\text{off}} \) at this stage of the algorithm.

Using \( S, S'_{\text{on}}, \) and \( S'_{\text{off}} \) is necessary because NTK-LRC is only useful as a comparative metric between two identically parametered-models. Since mutation implicitly modifies a model’s operations and therefore parameters, \( S \) and \( S'_{\text{on}} \) cannot be meaningfully compared. Using \( S'_{\text{on}} \) and \( S'_{\text{off}} \) remedies that issue, by creating two identically parametered models that mathematically represent a mutated and unmutated model respectively. As such, the NTK-LRC delta induced by the mutation can then be measured, which happens in line 7.

If this mutation produces a favorable NTK-LRC (that is, did not increase NTK and did not decrease LRC) it is added to the set of found good mutations \( \mathcal{E} \) in line 8. This whole

<table>
<thead>
<tr>
<th>Algorithm 1: NTK-LRC Mutation Selection</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>Input</strong>: Some model ( M ) and sample size ( n_{\text{good}} )</td>
</tr>
<tr>
<td><strong>Output</strong>: ( e^* )</td>
</tr>
<tr>
<td>Let ( s(x) ) refer to the VRAM size of some component ( x );</td>
</tr>
<tr>
<td>Let ( s_{\text{max}} ) be the maximum permitted VRAM allocation;</td>
</tr>
<tr>
<td>1 Create ( S ), the slim model of ( M );</td>
</tr>
<tr>
<td>2 for edge ( e ) in ( M ), selected in random order do</td>
</tr>
<tr>
<td>3 Let ( S' ) be an exact copy of ( S );</td>
</tr>
<tr>
<td>4 Mutate ( e ) in ( S' ), thus creating two new edges ( e' ) and ( e'' );</td>
</tr>
<tr>
<td>5 Create an identical copy of ( S' ) and label the two ( S'<em>{\text{on}} ) and ( S'</em>{\text{off}} );</td>
</tr>
<tr>
<td>6 Within ( S'_{\text{off}} ), disconnect ( e' ) and ( e'' );</td>
</tr>
<tr>
<td>7 Compute NTK and LRC for ( S'<em>{\text{on}} ) and ( S'</em>{\text{off}} );</td>
</tr>
<tr>
<td>8 if ( \text{NTK}<em>{\text{on}} \leq \text{NTK}</em>{\text{off}} ) and ( \text{LRC}<em>{\text{on}} \geq \text{LRC}</em>{\text{off}} ) then</td>
</tr>
<tr>
<td>9 Add ( e ) into the set of good mutation candidates ( \mathcal{E} );</td>
</tr>
<tr>
<td>10 Let ( e^* ) be the candidate mutation with best joint NTK-LRC in ( \mathcal{E} );</td>
</tr>
<tr>
<td>11 if ( s(M) + 2s(e^*) &lt; s_{\text{max}} ) then</td>
</tr>
<tr>
<td>Output: ( e^* )</td>
</tr>
<tr>
<td><strong>Output</strong>: ( \emptyset )</td>
</tr>
</tbody>
</table>
process repeats until a satisfactory number of good mutations are found or all edges have been sampled, at which point the algorithm exits the edge sampling loop (line 9). If any good edges are found, then the one with the most favorable joint NTK-LRC is selected in line 10 as the final mutation choice. If this mutation is performable on $M$ within the available memory space (line 11), it is returned as the selected mutation. Otherwise, no edge is returned.

3.8. SpiderNet Algorithm

With the mutation guiding NTK-LRC metric described, all necessary components for the full SpiderNet algorithm are present. This is stated in Algorithm 2.

The algorithm first starts by initializing a minimum viable model $M$. Next, the model is trained and pruned, after which a number of mutations are performed (if possible), then the whole process repeats until the model has reached sufficient size. Then the full model trains and prunes freely until it is fully trained.

4. Experiment Designs

Given the above algorithm and metrics, a variety of experiments can be designed over CIFAR-10 to evaluate their various efficacies, the details of which are presented in this section.

4.1. Model and Training Configuration

Each SpiderNet model has a few configurable hyperparameters, and each model presented in this section adheres to the following design configuration:

- **Maximum Permitted VRAM Space**: 20, corresponding to max space of a NVidia 3090.
- **Initial Convolution Channels**: 64
- **Number of Reductions**: 2
- **Operation Set**: Same as DARTS: Identity, 3x3 Max Pool, 3x3 Average Pool, 3x3 Separable Convolution, 5x5 Separable Convolution, 3x3 Dilated Convolution, 5x5 Dilated Convolution

This is the advantage provided by SpiderNet’s large-cell structure and dynamically-evolving search space: just two hyperparameters dictate the search space, initial convolution channels and total reduction count. The other two, VRAM space and operational choice, are dictated by hardware/deployment constraints and the domain of the given task respectively. Compare this to something like DARTS, which requires additional choices of total cell count, reduction spacing, per-cell node choice, node degree, number of input nodes per cell, and number of skip-connections to subsequent cells to define the search space. SpiderNet handles all of these automatically, dynamically adapting to best fit the task at hand.

In terms of training, all models use the following hyperparameters:

- **Epochs**: 600
- **Batch Size**: 64
- **Learning Rate**: .01, cosine annealed to 0
- **Dropout**: 0.2
- **Data Augmentation**: Random Cropping, Horizontal Flipping, Cutout, Normalization

4.2. Evolution via Joint NTK-LRC Metric

Models are trained and pruned in four epoch cycles. After each cycle, NTK-LRC is used to find up to three beneficial mutations to perform on the model, provided there is available VRAM space. The model’s weights are reset, then the process repeats. 15 total cycles of this are performed.
after which the model trains and prunes as desired for 600 epochs.

4.3. Random Search and Ablation Studies

Each SpiderNet run is compared against a random search, to evaluate SpiderNet’s ability to find good networks within the search space as per Yu et al. [16]. In order to do this, it is first necessary to enumerate the different ‘non-random’ components of the algorithm. In the case of SpiderNet, there are three: inter-cycle pruning, guided mutation by metric, and intra-training pruning, from lines 1, 3, and 4 of Algorithm 2 respectively.

Replacing guided mutation with a random measure is simple; rather than choose the edges with optimal metric values for mutation, an edge is selected at random. Each random run then attempts to perform as many random mutations as was attempted during the compared SpiderNet run, typically 45 (15 cycles of three mutations each). Next, inter-cycle pruning is replaced by random operation deletion, where the number of operations deleted per cycle is taken from the per-cycle deadhead rates (Algorithm 2, line 2) of the comparison SpiderNet run. Finally, intra-train pruning is replaced via random operation deletion, deleting an equal number of operations as were deadheaded during the final model training of the comparison SpiderNet run. The latter two components can also be simply removed (that is, not performing pruning at all), as a simple ablation study.

Four combinations of these random/non-random/removed components were chosen to provide the widest gamut of comparisons, the specifics of which are shown in Table 1.

<table>
<thead>
<tr>
<th>Label</th>
<th>Mutation</th>
<th>Inter-cycle</th>
<th>Intra-train</th>
</tr>
</thead>
<tbody>
<tr>
<td>Random 1</td>
<td>Random</td>
<td>Random</td>
<td>Random</td>
</tr>
<tr>
<td>Random 2</td>
<td>Random</td>
<td>None</td>
<td>None</td>
</tr>
<tr>
<td>Random 3</td>
<td>Random</td>
<td>None</td>
<td>Pruners</td>
</tr>
<tr>
<td>Random 4</td>
<td>Random</td>
<td>Pruners</td>
<td></td>
</tr>
</tbody>
</table>

Table 1. Details of the chosen random experiments

These four were chosen for their specific comparative properties. Namely, Random 1 is the pure-random run that replaces each guided component of Spider with its random equivalent, thus serving as a purely random reference point against which to compare all experiments against. Random 2 is a pruning ablation study, which in comparison to Random 1 will inform us of the effects of the random pruning performed. Random 3 evaluates how well guided-pruning can mesh with random mutation. This will be useful as direct comparison against a guided-run, to determine how much of its performance stems from the guided-mutation versus guided intra-train pruning. Finally, Random 4 evaluates how random pruning during the mutation stage (which increases the amount of space available for mutations, therefore leading to an increased number of mutations as more of the 45 opportunities will be successful) might benefit final performance.

5. Results

Table 2 compares the results of SpiderNet against the four random methods. Next, Table 3 shows a comparison between SpiderNet and other similar NAS methods from the literature. Since the whole point of SpiderNet is to minimize configurational tuning, reported is the first run of each configuration.

<table>
<thead>
<tr>
<th>Growth Strategy</th>
<th>Allotted VRAM</th>
<th>Test Acc.</th>
<th>Search Time</th>
<th>Total Time</th>
<th>Params</th>
</tr>
</thead>
<tbody>
<tr>
<td>SpiderNet</td>
<td>20</td>
<td>97.13%</td>
<td>5.4</td>
<td>74.9</td>
<td>8.8</td>
</tr>
<tr>
<td>SpiderNet</td>
<td>6.5</td>
<td>96.78%</td>
<td>5.8</td>
<td>36.2</td>
<td>4.0</td>
</tr>
<tr>
<td>Random 1</td>
<td>20</td>
<td>96.96%</td>
<td>-</td>
<td>69.9</td>
<td>12.9</td>
</tr>
<tr>
<td>Random 2</td>
<td>20</td>
<td>97.03%</td>
<td>-</td>
<td>81.4</td>
<td>16.8</td>
</tr>
<tr>
<td>Random 3</td>
<td>9.5</td>
<td>96.78%</td>
<td>-</td>
<td>30.5</td>
<td>4.0</td>
</tr>
<tr>
<td>Random 4</td>
<td>20</td>
<td>96.97%</td>
<td>-</td>
<td>85.4</td>
<td>11.1</td>
</tr>
<tr>
<td>Random 4</td>
<td>20</td>
<td>96.81%</td>
<td>-</td>
<td>84.7</td>
<td>9.1</td>
</tr>
</tbody>
</table>

Table 2. Results of the six SpiderNet and random runs conducted. Total run time refers to the start-to-finish wall clock time from initialization to fully-trained model. Search time consists of the wall clock time of Algorithm 2 from start until Line 4. All search and run times reported in this section are given in GPU hours, max VRAM in gigabytes, and parameter counts in millions.

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>Test Acc.</th>
<th>Params</th>
<th>Search Hours</th>
</tr>
</thead>
<tbody>
<tr>
<td>NAS-Net [17]</td>
<td>95.53%</td>
<td>7.1</td>
<td>37.755</td>
</tr>
<tr>
<td>ENAS Micro-search [10]</td>
<td>96.13%</td>
<td>38.0</td>
<td>7.7</td>
</tr>
<tr>
<td>ENAS Macro-search [10]</td>
<td>97.11%</td>
<td>4.6</td>
<td>14.4</td>
</tr>
<tr>
<td>DARTS First-order [6]</td>
<td>97.00 ± 0.14%</td>
<td>3.3</td>
<td>36</td>
</tr>
<tr>
<td>DARTS Second-order [6]</td>
<td>97.24 ± 0.09%</td>
<td>3.3</td>
<td>96</td>
</tr>
<tr>
<td>NAO (w/ weight sharing) [7]</td>
<td>97.07%</td>
<td>2.5</td>
<td>7</td>
</tr>
<tr>
<td>ProxylessNAS [1]</td>
<td>97.92%</td>
<td>5.7</td>
<td>N/A</td>
</tr>
<tr>
<td>PC-DARTS [15]</td>
<td>97.43%</td>
<td>3.6</td>
<td>2.4</td>
</tr>
</tbody>
</table>

Table 3. CIFAR-10 statistics for a variety of similar NAS models covered thus far.
6. Discussion

First, the overall ramifications of Tables 2 and 3 will be explored. Then, the table will be broken into specific subsets of runs whose comparisons provide interesting insights.

6.1. Overall Results

As compared to the literature, SpiderNet produces highly competitive models. While it is outperformed by algorithms like ProxylessNAS and PC-DARTS, it does so sans the extensive configurational choices of things like node count and cell structure required by these algorithms. In situations where the best-practices for these choices are unavailable, these algorithms would require extensive tuning not required by SpiderNet.

Additionally, the runs of lower VRAM specification demonstrate that the SpiderNet performs well under different constraints: apart from differing VRAM caps, these runs were configured identically. This demonstrates a versatility for performing NAS in different hardware domains, specifically one that comes with very little need for configuration tuning. While the configurations that performed well in the high-VRAM domain may not be optimal for low-VRAM domains (nor are they guaranteed to be optimal even in their intended range), SpiderNet finds very good models regardless. This is exciting from an “applied NAS” perspective, in that it minimizes the need to tune the algorithm to the desired domain. This means that SpiderNet has fulfilled the goals that were outlined in the introduction, wherein it was of interest to explore an algorithm free of the structural configuration tuning that weighs down fixed-structure algorithms like DARTS or ProxylessNAS.

In terms of the randomly mutated models, the non-pruned Random 2 produced the best performance, while the pure-random Random 1 produced the fastest model by runtime. This is an interesting result that ran contrary to expectations, which stemmed from prior differentiable pruning results [3, 4]. In these, differentiable pruning during training was shown to increase accuracy and decrease cost as compared to random pruning or no pruning at all.

From these expectations, the results of Random 2 and Random 3 are surprising; if the unpruned supernet of Random 2 performed so well, it might be expected that a similarly-grown supernet that could differentiably prune would again produce better accuracy more cheaply. However, Random 3 produced worse accuracy four hours slower than Random 2, although it did use 5.7M fewer parameters. The same could be said about the Random 1 and Random 4 pair; we might expect Random 4 to perform better due to differentiably pruning during training. However, Random 4 was significantly less accurate than Random 1 as well as around 15 hours slower in total runtime.

One possible explanation to these results is that all previous pruner experiments were performed over identical supernet, while these four random runs each have a different randomly generated structure. Given how drastically the structure of SpiderNet models can vary (see Figures 1, 2, and 3 in the Supplementary Materials for examples of this), the performance differences seen may be the result of the architecture’s implicit capabilities, more so than the effects of pruning. To fully decouple these effects, the random methods would all need to be performed over the same randomly-mutated model, that is, pick a random sequence of mutations that are then deterministically performed every time. However, this is a tangential line of experimentation to the main focus of the work, and thus was not prioritized.

6.2. Best NAS versus Best Random

The first analyzed pairing is the best SpiderNet run against the best random run, shown in Table 4:

<table>
<thead>
<tr>
<th>Growth</th>
<th>Allotted VRAM</th>
<th>Test Acc.</th>
<th>Search Time</th>
<th>Total Time</th>
<th>Params</th>
</tr>
</thead>
<tbody>
<tr>
<td>SpiderNet</td>
<td>20</td>
<td>97.13%</td>
<td>5.4</td>
<td>74.9</td>
<td>8.8</td>
</tr>
<tr>
<td>Random 2</td>
<td>20</td>
<td>97.03%</td>
<td>-</td>
<td>81.4</td>
<td>16.8</td>
</tr>
</tbody>
</table>

Table 4. The best random run (Random 2) versus the best SpiderNet run.

The most immediate conclusion that may be drawn from this is that there are only marginal accuracy differences between the SpiderNet and random runs, with the best SpiderNet run scoring only 0.10 percentage points better than the best random run. However, there is a large parameter difference between the two; the SpiderNet run used less than half the parameters of the random run. Additionally, the SpiderNet run from start to finish took around six hours less than the random run. These latter two points are particularly interesting; the whole random-search condemnation of NAS stems from random search providing a model of equivalent accuracy and parameter count “for free”. In any search space wherein every candidate model has an identical count of single-path edges (DARTS or ProxylessNAS, for example) this might be a valid assumption, as the single-operation-per-edge restriction means that every possible model contains an identical number of operations (equal to the total number of edges e in the model), meaning training times will only vary based on the computational complexity of the selected operations. Meanwhile, if operations (and thus per-edge parameter counts) are selected from a uniform distribution (such as in [5]), it is expected that the total number of parameters in any randomly generated model will also be roughly similar. To test this, 100,000 models were generated from a single-path, 57 edge search space via uniform operation selection. The distribution of parameter counts was highly normal (Shapiro-Wilk statistic 0.9996, p=4.1 × 10^{-12}), with a mean of 1.67M and standard deviation of 0.28M.
the Yu et al. [16] nor Li & Talwalkar [5] papers that started the random search discussion consider model training time a factor, and only the latter mentions parameter count, albeit only to point out the equivalent parameter counts of their random and NAS models.

Meanwhile, the SpiderNet search space is entirely unbounded; models with only three operations are valid members of the search space as are those with 300 or 3,000. Therefore, the time costs of training and parameter counts of any such model in the space varies wildly, and is a massively important factor in determining the quality of any given model. In this regard, the SpiderNet model’s speed is notable; not only does it provide an equivalently accurate (in fact, better) model than the best random-search model, this model is so parameter and time-efficient that its 5.4 hour search time handicap is more than doubly made up over the course of training. Additionally, the SpiderNet run only used around 15 GB of its maximum allotted 20 GB, while Random 2 used the full 20 GB, another efficiency not considered in random-search studies.

6.3. Best NAS versus Pure Random

Next is the comparison of the best SpiderNet run to the purely random run, shown in Table 5:

<table>
<thead>
<tr>
<th>Growth Strategy</th>
<th>Allotted VRAM</th>
<th>Test Acc.</th>
<th>Search Time</th>
<th>Total Time</th>
<th>Params</th>
</tr>
</thead>
<tbody>
<tr>
<td>SpiderNet</td>
<td>20</td>
<td>97.13%</td>
<td>5.4</td>
<td>74.9</td>
<td>8.8</td>
</tr>
<tr>
<td>Random 1</td>
<td>20</td>
<td>96.96%</td>
<td>-</td>
<td>69.9</td>
<td>12.9</td>
</tr>
</tbody>
</table>

Table 5. The fastest random run (Random 1) versus the best SpiderNet run.

The pure random run compares more favorably to the best SpiderNet run, with a faster training time and closer parameter count, at the expense of marginally worse accuracy than Random 2. However, it still has around 4 million more parameters than the SpiderNet run. These results show that random search is relatively competitive against SpiderNet, albeit with consistently worse parameter efficiency and accuracy.

6.4. Constricted NAS versus Constricted Random

<table>
<thead>
<tr>
<th>Growth Strategy</th>
<th>Allotted VRAM</th>
<th>Test Acc.</th>
<th>Search Time</th>
<th>Total Time</th>
<th>Params</th>
</tr>
</thead>
<tbody>
<tr>
<td>SpiderNet</td>
<td>6.5</td>
<td>96.78%</td>
<td>5.8</td>
<td>36.2</td>
<td>4.0</td>
</tr>
<tr>
<td>Random 2</td>
<td>9.5</td>
<td>96.78%</td>
<td>-</td>
<td>30.5</td>
<td>4.0</td>
</tr>
</tbody>
</table>

Table 6. The constricted random run versus the constricted SpiderNet run.

The runs in Table 6 comprise a set of runs where overall, top-end accuracy was less prioritized in favor of parameter-efficiency (the SpiderNet and Random 2 runs). These two perform nearly identically, with SpiderNet taking 5.7 hours longer due to its search time costs but requiring 3 GB less VRAM.

7. Conclusion

SpiderNet stands as a proof-of-concept of a minimally-configured NAS algorithm, one that aims to minimize the number of manual choices needed to produce excellent networks. By operating within an unbound, infinite search space that is specified by just two parameters (reduction count and initial convolution scale), SpiderNet has the capability to incorporate previously manual design parameters like per-cell node count automatically into its search space, drastically reducing the available tuning parameters required to produce good models.

Another point of interest amongst our results is the triangular mutation’s ability to produce very high quality networks from minimal starting conditions. In this sense, this work is comparable to that of Real et al. [12]’s Regularized Evolution, which likewise grew models from minimal conditions to state-of-the-art capabilities. Indeed, the whole concept of “minimal initial conditions” is directly inspired by the “trivial initial conditions” described in Real et al. [12]. However, what is interesting about the triangular mutation is that it can do this so quickly and consistently, without much need to carefully consider where it should occur. The random results show this clearly, as although they are all slightly worse in various metrics than the SpiderNet runs, they are still comparable in performance to various other NAS approaches.

7.1. The False Equivalence of Random Search

This strong accuracy performance of the randomly grown models does not invalidate the results of SpiderNet, in fact, the opposite. Crucially, it revealed that comparison to random search in these idiosyncratic search spaces can create a false-equivalency. If the entire end-to-end process of finding and training a model takes as long or longer for a random algorithm than NAS, then any accuracy similarities between the two become significantly less relevant. Section 6.2 make this very clear; spending just five hours searching is rewarded by increased accuracy, lowered VRAM, and halved parameter counts, and then the search cost is entirely refunded and more by the increased training speed. This completely erodes any argument that random search’s comparable accuracy detracts from NAS; such an argument entirely ignores the three other dimensions of parameter, VRAM, and time efficiency by which SpiderNet and NAS show clear benefits.
References


