





## **Overview**

In this paper we study the problem of **automatically generating polynomial solvers** for minimal problems. The main contribution is a new method for **finding small elimination templates** by making use of the syzygies (i.e. the polynomial relations) that exist between the original equations. Using these syzygies we can essentially parameterize the set of possible elimination templates. We evaluate our method on a wide variety of problems from geometric computer vision and show improvement compared to both handcrafted and automatically generated solvers. Furthermore we apply our method on two previously unsolved relative orientation problems.

### Background

- Systems of polynomial equations occur in many geometric vision problems.
- an eigenvalue problem

 $lpha(oldsymbol{x})$ 

 $\alpha(\boldsymbol{x})b_i(\boldsymbol{x})$ 

Solutions to (1) are then found by eigenvalue decomposition of the matrix M. • To find the action matrix M a so-called **elimination template** is used

$$CX = \mathbf{0} \tag{4}$$

where C is a constant matrix depending on the data and X is a vector of monomials.

- Selecting which monomials to multiply the equations with to form C is difficult.
- template is sufficient.

### Automatic Generator

We have implemented our approach in an automatic generator similar to that of Kukelova et al.[33].

Code is available at http://www.maths.lth.se/~viktorl/

#### References

[33] Z. Kukelova, M. Bujnak, and T. Pajdla. Automatic generator of minimal problem solvers. In *European* Conference on Computer Vision, pages 302–315. Springer, 2008.

# Efficient Solvers for Minimal Problems by Syzygy-Based Reduction

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$$\begin{cases} f_1(\boldsymbol{x}) = 0 \\ \vdots \\ f_m(\boldsymbol{x}) = 0 \end{cases}$$
(1)

• Most common method (in Computer Vision) is the **Action Matrix** method which transforms the problem to

$$\begin{bmatrix} b_1(\boldsymbol{x}) \\ \vdots \\ b_K(\boldsymbol{x}) \end{bmatrix} = \begin{bmatrix} M \end{bmatrix} \begin{bmatrix} b_1(\boldsymbol{x}) \\ \vdots \\ b_K(\boldsymbol{x}) \end{bmatrix}$$
(2)
$$E - \sum_{i} m_{ij} b_j(\boldsymbol{x}) = \sum_{i} h_{ij}(\boldsymbol{x}) f_j(\boldsymbol{x})$$
(3)



• Each row corresponds to an equation of the type  $\boldsymbol{x}^{\beta}f_{i}(\boldsymbol{x}) = 0$  where  $\boldsymbol{x}^{\beta}$  is some monomial.

• Forming linear combinations of the rows in C allows us to find the polynomials in (3).

• One approach is by Kukelova et al. [33] where the elimination template is found by an iterative approach. The method alternates between expanding the set of equations and performing gaussian elimination to check if the





#### Elimination templates

- However for a given problem the  $h_{ij}$  are **not unique**.

$$\operatorname{Syz}(f_1, \dots, f_m) = \left\{ \boldsymbol{s} = (s_1, \dots, s_m) \in \mathbb{K}[X]^m \mid \sum_k s_k(\boldsymbol{x}) f_k(\boldsymbol{x}) = 0 \right\}$$
(5)

elimination template.

$$\sum_{j}$$

- Compute normal form  $\overline{h_i}^{G_S}$  w.r.t. the Gröbner basis. Degree as a proxy for having few monomials.

#### Finding elimination templates

In this paper we propose the following approach for finding the elimination template for a given problem. 1. Generate an instance of the problem with coefficients in  $\mathbb{Z}_p$ 

- 3. Form the polynomials

and find coefficients  $h_{ij}$  satisfying (3) by division with G.



• We want to **minimize the number of monomials** (since these correspond to rows in C). • This ambiguity is characterized by the (first) syzygy module of  $(f_1, \ldots, f_m)$ .

So for any  $h_i = (h_{i1}, \ldots, h_{im})$  which satisfies (3) we can add any element in  $Syz(f_1, \ldots, f_m)$  and get another

$$h_{ij}f_j = \sum_j (h_{ij} + s_j)f_j = \sum_j \tilde{h}_{ij}f_j \quad \forall \boldsymbol{s} \in \operatorname{Syz}(f_1, \dots, f_m)$$
(6)

• Choosing the optimal  $\mathbf{s} \in \text{Syz}(f_1, \ldots, f_m)$  is a difficult problem.

• We propose a simple heuristic for simplifying the coefficients  $h_i$ .

Compute Gröbner basis  $G_S$  for the module  $Syz(f_1, \ldots, f_m)$  w.r.t. GRevLex-TOP.

2. Compute a Gröbner basis  $G = \{g_k\}$  and take  $\{b_1, \ldots, b_k\}$  as the standard monomials.

$$\alpha b_i - \overline{\alpha b_i}^G = \alpha b_i - \sum_j m_{ij} b_j \tag{7}$$

4. Compute Gröbner basis  $G_S$  for  $\operatorname{Syz}(f_1, \ldots, f_m)$  and form  $\tilde{\boldsymbol{h}}_i = \overline{\boldsymbol{h}}_i^{G_S}$ 5. Build template by multiplying the equations with the monomials from  $h_{ij}$ .

Each of these steps can be performed efficiently in a few lines of Macaulay2 code.

#### **Experimental evaluation**





#### Problem

Rel. pose 5pt Rel. pose 8pt one-TDOA offset rank Rel. pose + one fP3.5P + focalRel. pose + constRel. pose + rad. Rel. pose 6pt ones TDOA offset rank Rolling shutter po Generalized P4P Stitching + constTDOA offset rank TDOA offset rank Generalized rel. p Optimal PnP Triangulation from Optimal PnP (Ca P4P + focal + racRel. pose + rad. Rel. pose + 2 rad Rel. pose 7pt one-Weak PnP Weak PnP (2x2 sRolling shutter Re Optimal pose w d Rel. pose w dir. Rel. pose w dir. Abs. pose quivers  $L_2$  3 view triangu Rel. pose w angle Refractive P5P TDOA offset rank Optimal PnP Optimal PnP (usi Optimal pose w d Optimal PnP (qua Refractive P6P + Rel. pose + constDual-Receiver TI Optimal PnP (rot  $L_2$  3 view triangu



Histograms of residual errors for 5,000 runs – from left to right – image stitching with unknown focal length and radial distortion [8, 39], the optimal PnP-method of Hesch et al. [21] and the optimal PnP-method of Zheng et al. [54].

	Original		Proposed generator	
	Author	template size	no reduction step	with reduction step
	Stewénius et al.[47]	${f 10 imes 20}$	10 imes 20	f 10 imes f 20
e-sided rad. dist.	Kuang et al.[30]	$12 \times 24$	f 11 imes f 20	f 11 imes f 20
k 2, 7,4 pts	Kuang et al. $[28]$	${f 20 imes 15}$	${f 20 imes 15}$	${f 20 imes 15}$
focal 6pt	Bujnak et al. $[4]$ (*)	f 21 imes f 30	f 21 imes f 30	f 21 imes f 30
	Wu [52]	${f 20 imes 43}$	$24 \times 45$	$20 \times 44$
t. focal 6pt	Kukelova et al. $[33]$ (*)	<b>31</b> imes <b>46</b>	$31 \times 50$	$31 \times 50$
dist. 8pt	Kukelova et al. $[33]$ (*)	$32 \times 48$	<b>31</b> imes <b>49</b>	$32 \times 50$
es-sided rad. dist.	Kuang et al.[30]	$48 \times 70$	${f 34 imes 60}$	${f 34 imes 60}$
k 2, 5,6 pts	Kuang et al. $[28]$	$105 \times 83$	$105 \times 83$	f 40 imes f 42
ose	Saurer et al. $[44]$ (*)	$48 \times 56$	$50 \times 55$	f 47 imes f 55
+ scale	Ventura et al.[51] (*)	$48 \times 56$	$50 \times 55$	f 47 imes f 55
f. focal $+$ rad. dist. 3pt	Naroditsky et al. $[39]$	$54 \times 77$	$96 \times 108$	f 48 imes 66
k 3, 9,5 pts	Kuang et al. $[28]$	${f 70 imes 31}$	${f 70 imes 31}$	${f 70 imes 31}$
k 3, 7,6 pts	Kuang et al. $[28]$	$255 \times 157$	$255 \times 157$	f 75 imes 57
pose 6pt	Stewénius et al.[48]	$f 60 imes 120^{\ddagger}$	$135 \times 164$	$99 \times 163$
	Hesch et al. $[21]$	$120 \times 120$	$93 \times 116$	f 88 imes 115
m satellite im.	Zheng et al. $[53]$ (*)	$93 \times 120$	$93 \times 116$	f 88 imes 115
ayley)	Nakano [38] (*)	$124 \times 164$	$186 \times 161$	118 imes158
nd. dist.	Bujnak et al. $[5]$ (*)	$136 \times 152$	${f 140 imes 144}$	$140 \times 156$
dist. 6pt	Kukelova et al. $[33]$ (*)	$238 \times 290$	$223 \times 290$	f 154 imes 210
l. dist. 9pt	Kukelova et al. $[33]$ (*)	$179 \times 203$	$355 \times 298$	f 165 imes 200
e-sided focal $+$ rad. dist.	Kuang et al.[30]	$200 \times 231$	$249 \times 214$	f 185 imes 204
	Larsson et al. $[35]$	$234 \times 276$	$568 \times 498^{\dagger}$	$f 189 imes 232$ $^{\dagger}$
sym)	Larsson et al.[35]	$104 \times 90$	$83 \times 90^{\dagger}$	$f 49 imes f 59^\dagger$
6P	Albl et al. $[2]$ (*)	196 imes 216	$222 \times 230$	$204 \times 224$
lir 4pt	Svärm et al. [49]	$280 \times 252$	$371 \times 351$	f 203 imes 239
3pt	Saurer et al. $[45]$ (*)	$411 \times 489$	$287 \times 324$	f 210 imes 255
3pt (using sym.)	_	_	$94 \times 111^{\dagger}$	$f 40  imes f 57^{\dagger}$
5	Kuang et al. [27]	$372 \times 386$	$420 \times 406$	f 217 imes 253
ulation (Relaxed)	Kukelova et al. $[34]$ (*)	$274 \times 305$	$399 \times 384$	$239\times290$
e 4pt	Li et al.[36] (*)	f 270 imes 290	$280 \times 304$	$266 \times 329$
_	Haner et al. $[16]$	$280 \times 399$	$410 \times 480$	${f 240 imes 324}$
k 3, 6,8 pts	Kuang et al. $[28]$	$1359 \times 754$	$1359 \times 754$	${f 356 imes 345}$
	Zheng et al. $[54]$ (*)	$575 \times 656$	$812 \times 704$	${f 521 imes 601}$
ing sym.)	Zheng et al. $[54]$ (*)	$348 \times 376$	$484 \times 408^{\dagger}$	$302\times342^{\dagger}$
lir 3pt	Svärm et al.[49]	$1,260 \times 1,278$	$918 \times 726$	${\bf 544 \times 592}$
aternion)	Nakano [38] (*)	$630 \times 710$	$958 \times 693$	f 604  imes f 684
- focal	Haner et al.[16]	$648 \times 917$	$2,196 \times 1,913^{\dagger}$	$636\times851^{\dagger}$
t. focal $+$ rad. dist. 7pt	Jiang et al. [23]	$886 \times 1.011$	$1,393 \times 1,237$	${f 581 imes 862}$
DOA 5pt	Burgess et al.[7]	$2,625 \times 2.352$	$850 \times 1,167$	f 455 imes 768
t. matrix)	Nakano [38] (*)	$1,936 \times 1,976$	$1,698 \times 1,153$	<b>1</b> , <b>102</b> imes <b>1</b> , <b>135</b>
ilation	Kukelova et al. $[34]$ (*)	$1,866 \times 1,975$	$2,647 \times 2,584$	$\mathbf{1,759\times 2,013}$

(\*) Original template constructed using [33]. If several elimination templates are used, the largest of these templates is reported. †: The problem contains variable-aligned symmetries [3, 31, 35] that was automatically found and removed by our generator. ‡: The original template doesn't generate the full Gröbner basis, and some additional operations on the template are performed.